

Writing Your Own Wireshark Packet Dissectors (ADVANCED)

March 31, 2008

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SHARKFEST '08

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March 31 - April 2, 2008

Advanced dissector writing

Techniques needed for protocols that are “complicated”:

- Fragment reassembly
- Decryption and decompression
- Conversations and per-packet data
- Request/response matching
- Cleaning up allocated data
- Expert info
- Taps

“Simple” protocols

Many protocols are “simple”

- One PDU per bearer protocol PDU
- No encryption, compression, etc.
- Each PDU can be interpreted independently

“Complicated” protocols

But not all protocols are simple

- Fragmentation and reassembly
- PDU body encoding
- PDUs that require context from previous PDUs
- etc.

“Data sources”

Not all data in a dissected PDU comes directly from current lower-level PDU

- Can come from other lower-level PDUs (reassembly)
- Can come from decompressed/decrypted/etc. data

“Data sources” (cont’d)

Solution: named data sources

- Name, used as tab name in hex dump pane
- “Real data” tvbuff, containing data
 - Created with `tvb_new_real_data`
- Data source created with `add_new_data_source`
 - Added to list of data sources for packet
 - `add_new_data_source(pinfo, tvbuff, name)`

Dissector initialization routines

Dissector might need to do initialization/cleanup work when capture is opened

- Initialize internal data structures
- Free data left over from previous capture
- `register_init_routine()` registers callback - callback has no arguments and no return value

Fragmentation and reassembly

Some higher-layer PDUs are built from multiple lower-layer PDUs

- IP fragmentation, IEEE 802.11 fragmentation, etc.
- Need to reassemble
- Lower-layer PDUs contain header plus payload
- Payloads of fragments are reassembled into higher-layer PDU
- Header of fragment indicates where it appears in higher-layer PDU

Reassembly with ID and offset

Some protocols identify fragments with PDU ID and byte offsets

- Examples: IPv4, IPv6
- ID identifies reassembled PDU (e.g., IP ID)
- Byte offset is offset within reassembled PDU of first byte of payload
- Last fragment is specially marked

Process payload

`fragment_add_check()` does “heavy lifting” of reassembly

- The first time this packet is seen:
 - Just returns NULL if fragment cut short by snaplen
 - Adds to reassembly based on `pinfo->src`, `pinfo->dst`, `id`
 - If all fragments found, saves as finished reassembly and returns `fragment_data *` for finished reassembly
 - Otherwise, returns NULL
- All times after that, looks for finished reassembly and returns `fragment_data *` for finished reassembly

Process payload (cont'd)

`fragment_add_check()` uses two `GHashTable` structures to keep track of fragments and reassemblies

- Initialize in dissector initialization routine
- Initialize first with `fragment_table_init()`
- Initialize second with `reassembled_table_init()`

Dissect reassembled payload

`process_reassembled_data()` does the “heavy lifting” to process a possibly reassembled PDU

- Checks whether reassembly done (`head != NULL`)
- If so, checks whether there’s more than one fragment
- If more than one fragment:
 - Creates new `tvbuff` for reassembled data
 - Adds a data source for it, with specified name
 - Adds items with information about fragments to protocol tree
- If only one fragment, creates subset `tvbuff` for payload
- If reassembly not done, returns `NULL`

Dissect reassembled payload (cont'd)

`process_reassembled_data()` needs `ett_` and `hf_` variables for subtrees and items it adds

- `fragment_items` structure specifies the variables
- Contains pointers to the variables

Reassembly with sequence # and offset

Some protocols identify fragments with PDU ID and sequence number

- Examples: TDS, TIPC
- ID identifies reassembled PDU
- Sequence number is ordinal number of fragment
 - 0-based or 1-based
- Last fragment specially marked

Process payload

Similar to fragment byte offsets

- `fragment_add_seq_check()`
- Takes sequence # rather than byte offset as argument
- Sequence # is 0-based
- For protocols with 1-based sequence #'s, subtract 1

Dissect reassembled payload

Same as for fragment byte offsets

- `process_reassembled_data()` hides the differences

Decryption/decompression

Some protocols encrypt or compress data in the PDU

- Must decrypt or decompress before processing
- Cannot decrypt/decompress in place
- Must generate new data array and tvbuff for data array
- Make a data source from tvbuff

Decrypt/decompress into new buffer

Allocate a buffer and decrypt/decompress into it:

```
guint8 *buff;  
tvbuff_t *new_tvb;  
int actual_size, captured_size;
```

```
actual_size = amount of decrypted/decompressed data;  
captured_size = amount of that data we can generate from captured data;  
buff = g_malloc(captured_size);  
decrypt/decompress into buff;
```

Set up new tvbuff for buffer

Allocate a new tvbuff for the buffer:

```
new_tvb = tvb_new_real_data(buff, captured_size, actual_size);
```

Arrange that the buffer be freed when the tvbuff is freed:

```
tvb_set_free_cb(new_tvb, g_free);
```

Set up new data source

Arrange that the new tvbuff be cleaned up when the original tvbuff is cleaned up:

```
tvb_set_child_real_data_tvbuff(tvb, new_tvb);
```

Add the new tvbuff as a data source, so it shows up as a tab in the hex dump pane:

```
add_new_data_source(pinfo, new_tvb, name);
```

Do dissection with the new tvbuff

Conversations

Mechanism for keeping track of “flows” (connections, etc.)

- Identified by endpoint addresses and port numbers
- Addresses are address structures
 - Address type and value as bytes
- Ports are numbers
 - Conversation has “port type” identifying type of port
 - TCP, UDP, SCTP, IPX (socket), etc.

Conversation state

Information about flow attached to conversation

- Can set dissector for conversation
 - Protocol like SDP can indicate “Set up TCP session with this protocol”
 - Setting dissector means future packets will be dissected properly
- Can attach data to a conversation
 - Contains state information needed to dissect packets in conversation

Creating conversations

Use the function `conversation_new()`

- Caller must check for existing conversation first
- Arguments:
 - Frame number of first frame
 - Endpoint addresses
 - Port type and endpoint ports
 - Options
- Returns `conversation_t * handle`

Finding conversations

Use the function `find_conversation()`

- Arguments:
 - Frame number of first frame
 - Endpoint addresses
 - Port type and endpoint ports
 - Options
- Returns `conversation_t * handle`

Wildcards

`conversation_new()` options allow “wildcarding” of addresses and ports

- Can create conversation with one “wildcard” address and/or port
 - For UDP traffic where reply can come from address or port different from request destination
 - For “future” conversation where both endpoints are not known yet
- Can specify `NO_ADDR2` and/or `NO_PORT2`

Wildcard matching

`find_conversation()` options control “wildcarding” of address and port

- Either fully specified or wildcard conversations can match
- Match with fewest wildcards wins
- Can specify `NO_ADDR_B` and/or `NO_PORT_B`
- Wildcard address or port can match *either* endpoint

Completing wildcard conversations

Matching can cause wildcards to be filled in

- Filling in is used for “future” connection-oriented conversations
- Protocol such as SDP can indicate “Set up TCP session with this protocol”
- Perhaps only one endpoint is known
 - Receiver of SDP message will connect to that endpoint
 - Full endpoint from which it connects is unknown
- Once connection is made, wildcarded endpoint is filled in by `find_conversation()`

Conversation dissector

Set with `conversation_set_dissector()`

- Takes `conversation_t *` and `dissector handle` as arguments
- Works for TCP, UDP, Datagram Congestion Control Protocol (DCCP), AppleTalk Transaction Protocol
- Takes precedence over heuristics and port matches

Conversation data

Each protocol can attach data to a conversation

- Data is opaque - not interpreted by conversation code
- Allocating and freeing is the dissector's responsibility
- Check whether data already exists, then add data if it does not exist
- Data can be changed as you dissect packets in the conversation
 - Data should be changed only on first pass through packets - `pinfo->fd->flags.visited false`

Adding conversation data

Use the function `conversation_add_proto_data()`

- Arguments:
 - `conversation_t * handle`
 - Protocol number for protocol
 - `void *` pointing to data
- No return value

Finding conversation data

Use the function `conversation_get_proto_data()`

- Arguments:
 - `conversation_t * handle`
 - Protocol number for protocol
- Returns `void *` pointing to data, or `NULL` if no data for protocol

Per-packet data

Data needed in order to dissect a particular packet correctly

- Might come from previous packets
- Might come from per-conversation data updated by previous packets

Adding per-packet data

Use the function `p_add_proto_data()`

- Arguments:
 - `frame_data * handle`
 - Protocol number for protocol
 - `void *` pointing to data
- No return value
- Check whether already present before adding
 - Data should be added only on first pass through packets

Finding per-packet data

Use the function `p_get_proto_data()`

- Arguments:
 - `frame_data * handle`
 - Protocol number for protocol
- Returns `void *` pointing to data, or `NULL` if no data for protocol

Request/response matching

Many protocols are request-response protocols

- Decoding response might require info from request
- User might want response to show frame # of request
- User might want time between request and response

Request/response matching table

If multiple requests in flight, protocol probably has request ID field

- Use `GHashTable` or `se_tree` with request ID as key
 - One table per conversation, not one global table!
- Store relevant information as value
 - Information needed to dissect response (e.g., request type)
 - Time stamp of request
 - Frame #s of request and response (0 means unknown)

Freeing allocated data

Data allocated privately by dissectors eventually needs to be freed

- If `se_` allocators used, freeing happens automatically
- Otherwise, need to free data in your initialization routine

Expert analysis

Log of “possibly interesting” behavior in a capture

- Allows users to get a summary of what they might want to look at
- Four severity levels:
 - Chat - interesting events in normal traffic flow, such as TCP SYN
 - Note - notable but not unusual events, such as HTTP 404
 - Warn - unusual events, such as a connection failure
 - Error - serious problem, such as a malformed packet

Expert information groups

The general type of condition an item describes

- Bad checksum
- Protocol sequence problem (discontinuous sequence numbers, retransmissions, etc.)
- Error response (gives error code)
- Request code (typically at Chat level)
- Undecoded
- Reassembly problem
- Malformed packet

Setting expert info

“Expert info” is a property of a protocol tree item

- Must have an item to which expert info is attached
- Added with `expert_add_info_format()`:
 - `packet_info * (pinfo)` pointer
 - `proto_item * pointer` to item
 - group
 - severity
 - `printf`-style format string and arguments

Taps and tap listeners

Mechanism for producing statistics, etc., from packets

- Dissectors provide taps
- Statistics routines provide tap listeners
- Tap listeners attach to taps

Taps

Taps can supply pre-digested data to listeners

- Register tap by name in `proto_register_routine`
- Queue packet for tap in dissection routine
 - Pass (non-auto!) data structure with pre-digested data, if tap supplies any

Tap listeners

Tap listeners process data supplied by taps

- Per-packet “packet” callback arguments are:
 - `pinfo`
 - Dissection information (including protocol tree)
 - Pre-digested information from dissector, if any
- “Display” callback called when accumulated information should be displayed or updated
- Tap listeners with UI cannot be shared between Wireshark and TShark - with an exception...

stats_tree taps

stats_tree tap does the UI work for you

- Displays a tree view
 - A list view is just a tree view with one level
- All you do is add nodes and update nodes
- Works with both Wireshark and TShark
- stats_tree tap can be a plugin (see stats_tree plugin)

Further information

- Everything Gerald mentioned in “Further information” in the previous session
 - (Except for “Next session”; no infinite looping here :-))
- `doc/README.request_response_tracking`
- `doc/README.tapping`
- `doc/README.stats_tree`

Q&A



Bonus material

Sample code



Process payload in reassembly

```
if (doing defragmentation && this is part of fragmented PDU) {
    head = fragment_add_check(tvb, payload_offset, pinfo, ID,
        fragment_table, reassembled_table, byte_offset,
        fragment_data_len, true_for_last_fragment);
    next_tvb = process_reassembled_data(tvb, offset, pinfo,
        "name", head, &frag_items, &update_col_info, tree);
} else {
    if (this is not part of fragmented PDU || this is the first fragment) {
        next_tvb = tvb_new_subset(tvb, payload_offset, -1, -1);
        if (part of fragmented PDU)
            pinfo->fragmented = TRUE;
        else
            pinfo->fragmented = FALSE;
    } else
        next_tvb = NULL;
}
```


Do dissection after reassembly

```
if (next_tvb == NULL) {
    /* Just show this as a fragment. */
    if (check_col(pinfo->cinfo, COL_INFO)) {
        col_add_fstr(pinfo->cinfo, COL_INFO, something to
            show this as a fragment);
    }
    if (head && head->reassembled_in != pinfo->fd->num) {
        if (check_col(pinfo->cinfo, COL_INFO)) {
            col_append_fstr(pinfo->cinfo, COL_INFO,
                " [Reassembled in #%u]", head->reassembled_in);
        }
    }
    call_dissector(data_handle, tvb_new_subset(tvb,
        payload_offset, -1, -1), pinfo, parent_tree);
    pinfo->fragmented = save_fragmented;
} else {
    hand next_tvb to the next dissector;
}
```